

SRBAC : Secure Role Base Access Control in Dynamic Cloud

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Abstract- The Attribute Base Encryption (ABE) provides the secure data storage and data sharing to data owners in cloud environments without any third party dependency. The former ABE schemes contain just one authority to be used for end user verification, which always generates a single-point bottleneck on both security as well as performance. Therefore, in existing too many multi-authority schemes are proposed, in which multiple authorities can maintain the privacy and verification of multiple attribute subsets. Even though, in many systems single-point bottleneck issues are not solved. In this proposed research work, from a system focus, a multi-authority verification scheme called Robust Access Control (RAAC) which is the part of Role Base Access Control (RBAC), the system proposed a secure data transmission in cloud environments. In RAAC, taking benefit of $(t$ out of n) threshold secret sharing, the master key and private keys can be shared amongst multiple authorities as well as a trusted party, and an authenticated user can generate his/her secret key by cooperating with any available authorities during the transaction. The ElGamal encryption provides the highest security from data with two types of cipher text policy and MK and PK with multiple authorities. Furthermore, by combining the conventional multi-authority scheme with RAAC, we design an approach that can verify multiple users in parallel and provide consistent data to end users. The experimental analysis also shows the proposed schemes are better than existing approaches.

Keywords- RBAC, ElGamal encryption scheme; secure user Revocation; Proxy Key Generation, Role Base Access Control (RBAC)

I. INTRODUCTION

Forward data security during the data transmission as well as storage in cloud environments is today's need for all applications. It moreover gives the authenticity and namelessness of the end user. Ring topology is the promising prospect to build an unknown and credible information sharing framework. It authorizes an information owner to their mystery validate his information which can be put into the cloud for capability or examination explanation. The proposed framework works in all cloud environments with multi-user verification approaches. The Trusted Third Party (TTP) and multiple Attribute Authorities (AA) provide the user verification that can remove the end user certificate verification process that can reduce the internal overhead of the system.

There are too many schemes that have been developed by the existing authors for cloud data security. Some Identity based, Key Base, CipherText based etc. each system having some drawbacks like collusion issues [3], bottleneck issues, SQL injection issues etc. Every scheme finds the new security approach for eliminating such drawbacks with new techniques. In this work, the system proposes a RAAC with CP-ABE access control plan, which eliminates the single-point queue on both security and execution. In this plan, numerous powers mutually deal with the entire property set however nobody has full control of a particular characteristic.

II. LITERATURE SURVEY

According to KaipingXue [1] propose a novel heterogeneous framework to remove the problem of single-point performance bottleneck and provide a more efficient access control scheme with an auditing mechanism. Our framework employs multiple attribute authorities to share the load of user legitimacy verification. Meanwhile, in our scheme, a CA (Central Authority) is introduced to generate secret keys for legitimacy verified users. Unlike other multi-authority access control schemes, each of the authorities in our scheme manages the whole attribute set individually. To enhance security, we also propose an auditing mechanism to detect which AA (Attribute Authority) has incorrectly or maliciously performed the legitimacy verification procedure.

Kan Yang and et. Al.[2], proposed a revocable multi-authority CP-ABE scheme, and apply it as the underlying techniques to design the data access control scheme. Our attribute revocation method can efficiently achieve both forward security and backward security. System also designs an expressive, efficient and revocable data access control scheme for multi-authority cloud storage systems, where there are multiple authorities co-exist and each authority is able to issue attributes independently.

The system [3] proposed a secure way for anti-collusion key distribution without any secure third party channels, and the users can securely get their private keys from group owner. Second, this method can propose fine-grained access control, any user in the group can use the source in the cloud and revoked users cannot access the cloud all over again after they are revoked. Thirdly, system can shield the scheme from collusion attack, that means that revoked users cannot get the actual data file even if they combine with the untrusted cloud. In this approach, by exploiting polynomial capability, framework can complete a safe client negotiation conspiracy. Finally, this plan can accomplish fine efficiency, which implies past clients necessitate not to refresh their revoked from the group.

According to [4] proposes the major of key-approach feature which is based on KP-ABE with reflection of non-monotonic access structures and with regular ciphertext size. System also proposes the first Key-Policy Attribute-based Encryption (KPABE) method allowing for non-granted access structures (i.e., that may contain negated attributes) and with constant ciphertext size. Towards achieving this goal, system first show that a certain class of identity based broadcast encryption schemes generically yields monotonic KPABE systems in the selective set model. System then describes a new efficient identity-based revocation mechanism that, when combined with a particular instantiation of our general monotonic construction, gives rise to the first truly expressive KP-ABE realization with constant-size ciphertext.

According to F. Zhang and K. Kim [5] proposed an ID-based ring signature approach, both approaches has defined base on bilinear pairings as well as Java pairing library. Also system analyzes their security and efficiency with different existing strategies. The Java Pairing library (JPBC) has used for data encryption and decryption purpose. Some user access control policies has design for end users that also enhance the privacy and anonymity of data owner.

In approach [6], propose the first Identity-based threshold ring signature approach that does not support to java pairings. It propose the first Identity -based threshold verifiable ring signature strategy. System also analyze that the secrecy of the actual signers is maintained even against the PK generator (PKG) of the Identity -based system. Finally system shows how to add identity collusion and other existing base different schemes. Due to the dissimilar levels of signer inscrutability they support, the system proposed in this paper actually form a suite of Identity -based threshold ring signature method which is related to many real-world systems with varied anonymity needs.

In [7], system first validates the security requirements of whole architecture, and after that adds to in the security architecture. System proposed AES 128 16 bit encryption approach for end to end user verification and data encryption/ decryption purpose.

According to Kan Yan [10], System proposed Ciphertext-Policy Attribute-based Encryption (CP-ABE) is a promising technique for access control of encrypted data. It requires a trusted authority manages all the attributes and distributes keys in the system. In cloud storage systems, there are

multiple authorities co-exist and each authority is able to issue attributes independently. However, existing CP-ABE schemes cannot be directly applied to data access control for multi-authority cloud storage systems, due to the inefficiency of decryption and revocation. In this paper, system propose DAC-MACS (Data Access Control for Multi-Authority Cloud Storage), an effective and secure data access control scheme with efficient decryption and revocation. Specifically, system construct a new multi-authority CP-ABE scheme with efficient decryption and also design an efficient attribute revocation method that can achieve both forward security and backward security

III. PROPOSED SYSTEM MODEL

The proposed system having 6 different objects, which work independently in different phases. Below we have shown all the entities

- 1) **Certificate Authority called as CA**
- 2) **Attribute Authorities called as AA's**
- 3) **Data Owner who upload the data**
- 4) **User who download the data**
- 5) **Trusted Middleware**
- 6) **Cloud Data Storage server**

A) *Implementation Details and Procedure:*

1. All Attribute Authorities has register with CA and get the unique Authority ID like $\{A1 \dots An\}$
2. All accessible users has also registered with CA
3. All users when access the data from cloud servers it will communicate with multiple AA's and Trusted Middleware and gets the private keys for data decryption.
4. All data owner collect the private keys from CA when upload any plane text, for decryption purpose.
5. After that data owner upload the encrypted text on to cloud server and same time distribute both keys to AA's and trusted middleware.
6. With the help of both keys user can download cipher data from cloud server.
 - o The system can achieve secure data sharing approach and access control. The RBAC is also proposed for data access and revocation has used for prevention for unauthenticated usage,
 - o Cost, secure in the sense that it can accomplish both backward security and forward security.

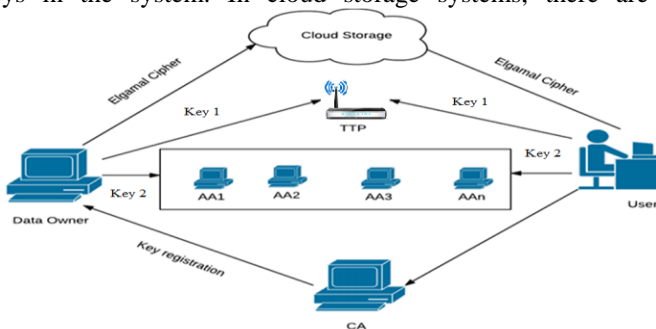


Fig.1: proposed system architecture

The above figure 1 shows the overall system architecture with all different entities. The master entity CA executes all necessary requirements of system. CA gives all four keys for data privacy, when data owner generates any requirements. With the help of those keys data owners generate encryption and distributed the master key and private key to middle ware entities. The system also focuses on data sharing approach like Role Base Access Control (RBAC). Data owner can share the owner data to different users and those users can access the same data from middle ware parties and multiple authorities. The verification has done from any available AA. The system also enhance the secure user revocation approach for data owner and proxy key generation against old keys. The overall scenario can eliminate the certificate verification process and bottleneck issues.

B) Mathematical Module:

The proposed system having six different module set
 $S = \{Dow, CA, TPA, CloudDB, AA's, User\}$

Data owner can upload the different text documents with respective key set

$$Dow = \{Dm1\langle k1 \dots kn \rangle, Dm2 \dots Dmn\langle k1 \dots kn \rangle\}$$

Generate a cipher data when upload

$$CData = \{Dm[i]\langle k1 \dots kn \rangle\}$$

TPA holds the all master keys which distributed by data owner

$$TPA = \{F1\langle mk \rangle, F2\langle mk \rangle \dots Fn\langle mk \rangle\}$$

AA's also holds all private keys of each file

$$AA's = \{F1\langle pk \rangle, F2\langle pk \rangle \dots Fn\langle pk \rangle\}$$

Data owner can share the specific file to particular user

$$ShareGroup = \{Uid\langle F, PK, MK \rangle[1], \dots Uid\langle F, PK, MK \rangle[n]\}$$

When user request for any file for cloud server, cloud server

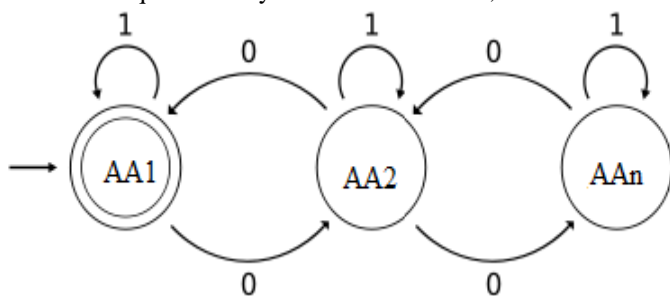


Fig.2 : State for each authority verification

Any t(n) return 1 then it will provide the private key otherwise state has change to another authority.

$$M = (Q, \Sigma, \delta, q_0, F) \text{ where}$$

$$Q = \{S_1, S_2\},$$

$$\Sigma = \{0, 1\},$$

$$q_0 = S_1,$$

$$F = \{S_1\}, \text{ and}$$

The proposed state diagram shows the how users request processed by middle ware authority. When request has generated from end user it first receives by AA. If the AA1 is already acquired by some other process or busy then it will forward to another AA. This process works like recursively or base on round robin approach, when any t authority gets free it will return the keys to authenticated user.

4.3 Algorithms:

1. Elgamal Encryption scheme

Key Generation phase

Input : Plain text as text data d.

Output: a,b,p,g all which contain public key, master key and private key

Step 1: Initialize the random message from user as d. (it should be any kind of text data).

Step 2: initialize a,b,p,g for private key purpose.

Step 3: generate P as randomly base on bit length of d. so, $Ans[] = \text{GetRandomP}(d.\text{getbyte}().\text{bitlength}$ base on probable prime no.

Step 4: $p = Ans[0]$

$$g = Ans[1]$$

Step 5: Generate a using P

$$a = \text{RandomA}(p)$$

its calculate like $p.\text{bitLength}()-1, \text{Random}.$

Step 6: Calculate $b = \text{calculateb}(g, a, p);$

$$\text{so, } b = g.\text{modPow}(a, p);$$

Step 7: Key generation done

Encryption

Input : Text data d,p,b,g

Output cipher as C1, and C2.

initialize $\text{BigInteger} [] \text{ rtn} = \{\text{null}, \text{null}\};$

$\text{message} = d.\text{getBytes}();$

$[] \text{ result} = \text{ElGamal}.\text{encrypt}(\text{message}, p, b, g);$

$[] \text{ rtn} = \{\text{null}, \text{null}\};$

$k = \text{ElGamal}.\text{getRandomk}(p);$

$C1 = g.\text{modPow}(k, p);$

$C2 = m.\text{multiply}(b.\text{modPow}(k, p)).\text{mod}(p);$

Decryption

Input : input c1 and c2 as cipher a and p as private keys

Output: Plain text d.

Step 1: $m = C2.\text{multiply}(C1.\text{modPow}(a.\text{negate}(), p)).\text{mod}(p);$

Step 2: return m.

2. SQL Injection and prevention algorithm for Database Security

1: Procedure $\text{QueryVerify}(\text{PlainQuery}, \text{PatternList}[])$

INPUT: $\text{PlainQuery} = \text{User input Query}, \text{Threshold } T$

$\text{PatternList}[] = \text{user define Pattern List with } m \text{ AnomalyPattern}$

2: For $j = 0$ to m do

3: If $(\text{Function}(\text{PlainQuery}, \text{String.Length}(\text{PlainQuery}), \text{PatternList}[j][0]) = 0)$ then

- 4: Calcsimilarity score of input query
- 5: If (Score Value Anomaly = T)
- 6 : Block query or stop query execution
- 7 : else execute on database

IV. SOFTWARE REQUIRMENTS

1. **System interfaces:** Ubuntu Operating System
2. **User interfaces:** User interface using Jsp and Servlet

3. Hardware interfaces

Processor :- Intel R-Core i3 2.7 or above

Memory :- 4GB or above

Hard Disk :- 500 GB

4. Software interfaces:

Front End: Jdk 1.7.0, Eclipse

IE 7.0/above

Back-End: Mysql 5.1.

5. Communications interfaces

System will use HTTP as well as SMTP and SOAP protocol for establishing connection and transmitting data over the network.

6. **Services:** Amazon EC2 as Public cloud Environemnt

V. EXPERIMENTAL RESULT

For the system performance evaluation, calculate the matrices for accuracy. The system is executed on java 3-tier architecture framework with INTEL 2.8 GHz i3 processor and

4 GB RAM with public cloud Amazon EC2 consol. For the system evaluation we create 2 machines on physical environment with Wi-Fi and 10 VM with Amazon EC2 as public cloud environment. After implementing some part of system we got system performance on reasonable level. The below table 1 shows the proposed Elagamal algorithm performance for user plain text conversion as well encryption decryption.

Data Size in MB	Encryption time (Milliseconds)		Decryption time (Milliseconds)	
	Existing	Proposed	Existing	Proposed
5	595	515	724	612
10	1120	1026	1132	1033
15	1680	1547	1687	1556
20	2260	2064	2231	2033

Table 1: System performance (Estimated)

In second experimentation system show the user verification time with different approaches. In current system we consider as four different authorities for runtime verification. The below Fig. 3 shows the performance measures using different parameters with some existing approaches.

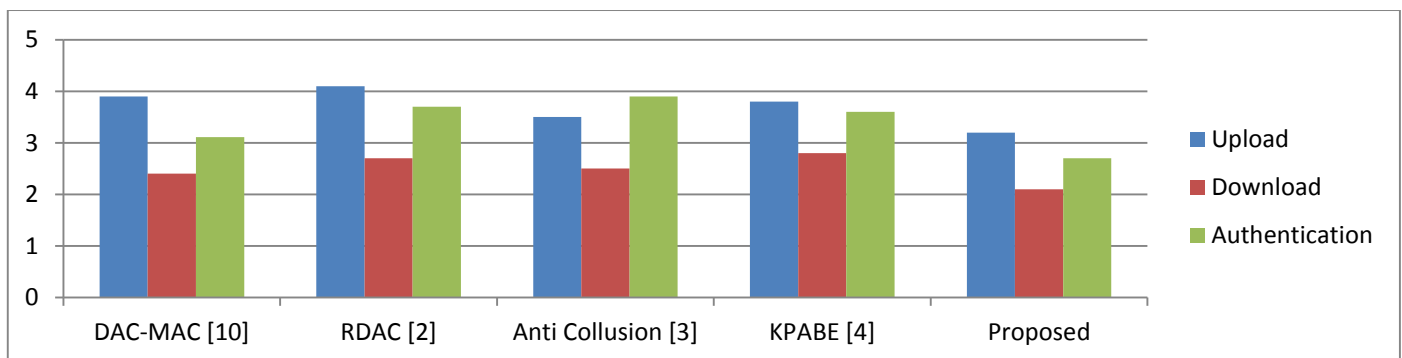


Fig.3: System Performance Measures proposed vs Existing approaches

VI. CONCLUSION

In this work system propose a secure Role Base Access Control (RBAC) data sharing scheme for untrusted environment in the cloud. In our scheme, the users can securely get their master and private keys from middleware authorities, CA provide and secure communication between multi parties. Also, our scheme is able to provide the secure revocation for untrusted user. The proxy key generation has also proposed in this work. When data owner revokes any specific end user system automatically expired the existing keys and generates new keys for all shared users. The system can achieve highest level security as well as privacy through such approaches.

VII. FUTURE WORK

The current architecture is very efficient for security purpose, but sometime its utilized multiple resources. When the such

system allocate multiple resources it will generate a lot of dependencies. For the next updation we can focus on minimum resource utilization with system flexibility like power, VM's, network, memory etc.

VIII. REFERENCES

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