



School of Engineering

Discrete Structures CS 2212 (Fall 2020)

15 – Induction and Recurrence

Sequence: A special type of function in which the domain is a consecutive set of integers.

| index \rightarrow | k | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|---------------------|------|---|----|----|----|----|----|-----|-----|-----|-----|-----|
| term→ | g(k) | 6 | 13 | 23 | 33 | 54 | 83 | 118 | 156 | 210 | 282 | 350 |

 $g(k) = g_k$

Finite sequence: A sequence with a finite domain.

Infinite sequence: A sequence with an infinite domain.

- Initial / Final index.
- Initial / Final term.

A sequence can be specified by an **explicit formula** showing how the value of term a_k depends on k.

Example:

| k | | | | | | | | | | | |
|-------|---|---|---|---|----|----|----|-----|-----|-----|------|
| a_k | 1 | 2 | 4 | 8 | 16 | 32 | 64 | 128 | 256 | 512 | 1024 |

Here, $a_k = 2^k$ for $k \ge 1$.

Increasing: For *every* two consecutive indices, *k* and *k* + 1, in the domain $a_k < a_{k+1}$.

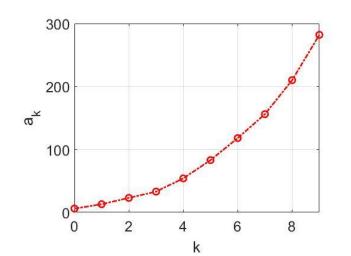
Non-decreasing: For *every* two consecutive indices, k and k + 1, in the domain, $a_k \le a_{k+1}$.

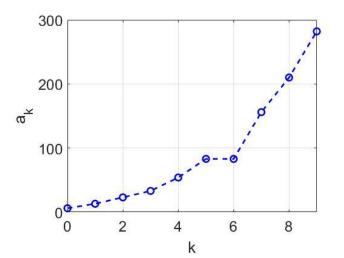
| k | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|-------|---|----|----|----|----|----|-----|-----|-----|-----|
| a_k | 6 | 13 | 23 | 33 | 54 | 83 | 118 | 156 | 210 | 282 |

| k | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|-------|---|----|----|----|----|----|----|-----|-----|-----|
| a_k | 6 | 13 | 23 | 33 | 54 | 83 | 83 | 156 | 210 | 282 |

Increasing

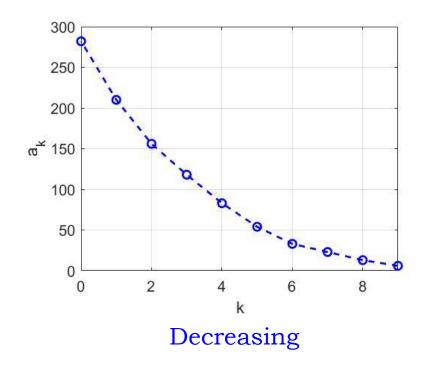


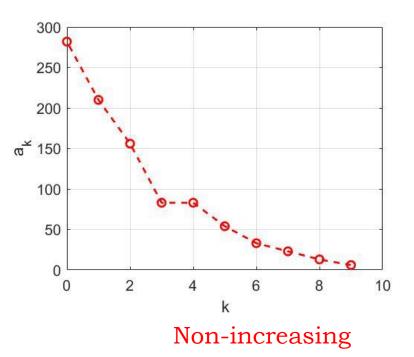


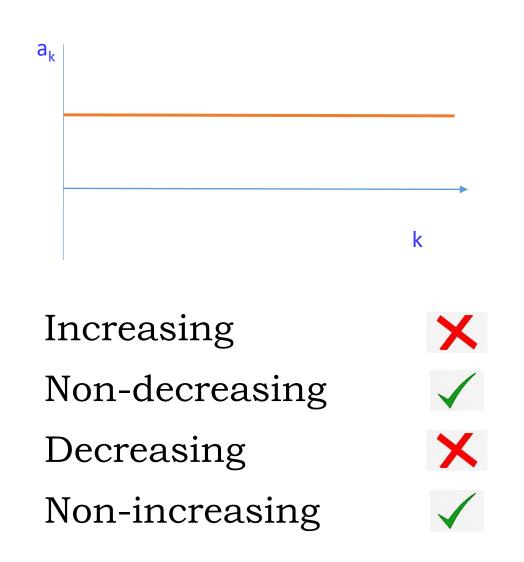


Decreasing: For *every* two consecutive indices, k and k + 1, in the domain $a_k > a_{k+1}$.

Non-increasing: For *every* two consecutive indices, k and k + 1, in the domain, $a_k \ge a_{k+1}$.







| a ₀ | a ₁ | a ₂ | a ₃ | a ₄ | ••• | |
|-----------------------|-----------------------|-----------------------|-----------------------|-----------------------|-----|--|
| | | | | | | |

| a ₀ | a ₁ | a ₂ | a ₃ | a ₄ | ••• | |
|-----------------------|-----------------------|-----------------------|-----------------------|-----------------------|-----|--|
| a_0 | | | | | | |

| a ₀ | a 1 | a ₂ | a ₃ | a ₄ | ••• | |
|----------------|--------------------------------------|-----------------------|----------------|-----------------------|-----|--|
| a_0 | (<i>a</i> ₀ + <i>d</i>) | | | | | |

| a ₀ | a ₁ | a ₂ | a ₃ | a ₄ | ••• | |
|----------------|------------------------------|------------------------------|-----------------------|-----------------------|-----|--|
| a_0 | (a ₀ + d) | (a ₁ + d) | | | | |

| a ₀ | a ₁ | a ₂ | a ₃ | a ₄ | ••• | |
|-----------------------|-----------------------|------------------------------|-----------------------|-----------------------|-----|--|
| a_0 | $(a_0 + d)$ | (a ₁ + d) | $(a_2 + d)$ | | | |

| a ₀ | a ₁ | a_2 | a ₃ | a ₄ | ••• | |
|-----------------------|-----------------------|-------------|-----------------------|------------------------------|-----|--|
| a_0 | $(a_0 + d)$ | $(a_1 + d)$ | $(a_2 + d)$ | (a ₃ + d) | ••• | |

- Here, $a_k a_{k-1} = d$, for all k > 0.
- Note that $a_k = a_0 + kd$
- Observe a "linear" growth.

| a ₀ | a ₁ | a_2 | a_3 | a ₄ | ••• |
|-----------------------|-----------------------|-------------|------------------------------|------------------------------|-----|
| a_0 | $(a_0 + d)$ | $(a_1 + d)$ | (a ₂ + d) | (a ₃ + d) | ••• |

Example:

- Suppose a person inherits a collection of 500 baseball cards and decides to continue growing the collection at a rate of 10 additional cards each week.
- A sequence of cards at the end of each week is an arithmetic sequence.

Geometric Sequences

A sequence where each term after the initial term is found by taking the previous term and **multiplying** by a fixed number called the **common ratio**.

| a ₀ | a ₁ | a 2 | a ₃ | ••• | a _k |
|-----------------------|-----------------------|----------------|-----------------------|-----|-----------------------|
| a_0 | $r \times a_0$ | $r \times a_1$ | $r \times a_2$ | ••• | $r \times a_{k-1}$ |

- Here, $\frac{a_k}{a_{k-1}} = r$, for all k.
- Note that $a_k = r^k a_0$
- Observe a "non-linear" growth.

Example: Money in a bank account earning a fixed rate of interest can be expressed as a geometric sequence.

Recurrence Relation: A rule that defines a term a_n as a *function of previous terms* in the sequence is called a recurrence relation.

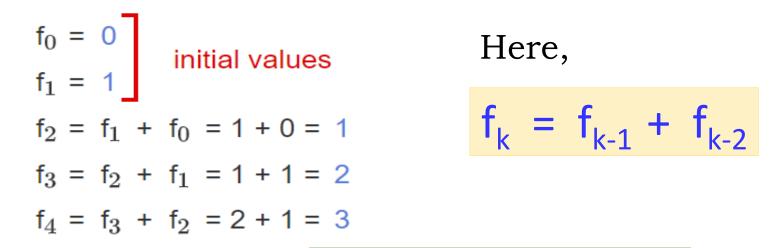
Arithmetic sequence

a₀ = a (initial value) a_n = d + a_{n-1} for n ≥ 1 (recurrence relation) Initial value = a. Common difference = d.

Geometric sequence

a₀ = a (initial value) a_n = r · a_{n-1} for n ≥ 1 (recurrence relation) Initial value = a. Common ratio = r.

Sometimes, we need multiple previous values to obtain the current term in the sequence. For instance,



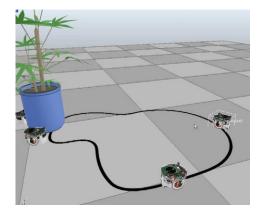
Such a sequence is known as Fibonacci sequence.

- Many applications: coding, optimization, algorithms analysis, arts etc.
- Golden ratio
- Miles to Kilometers (fun application)

Other applications of recurrence relations include modelling of (discrete time) **dynamical systems**.

Difference equations are used to model dynamics of such systems.

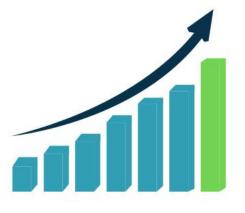
Simplest difference equation has the form.



Robot dynamics



Population growth



Economic models

Example: Describe the scenario below as a *recurrence relation*.

- The salary of an employee in year n is s_n
- The employee receives a \$1000 bonus in December.
- The salary for the subsequent year is 5% more than the income received by the employee in the previous year (including salary and bonus).



Answer:

- The income earned in year n 1 is $(s_{n-1} + 1000)$.
- Thus, the increase in salary is 5% of $(s_{n-1} + 1000)$.
- If a quantity increases by 5%, then the quantity becomes 1.05 times larger.
- Therefore the salary in year *n* can be described by the recurrence relation $1.05(s_{n-1} + 1000)$.

Sequences and Summations

Summation notation is used to express the sum of terms in a numerical sequence

Summation form:
$$\sum_{i=1}^{4} i^2$$
Expanded form: $1^2 + 2^2 + 3^2 + 4^2$

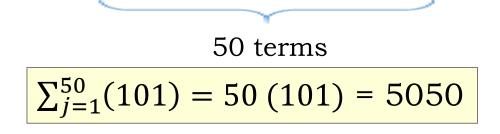
Always use parentheses as necessary to indicate which terms are included in the summation as in $\sum_{i=1}^{n} (i+1)$

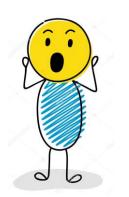
- Sometimes, it helps to "manipulate" the summation expression and make things easy, for instance, to compute the exact value of the expression, or to simplify the form of the expression.
- We will take a look at some tricks and tips to do so.
- Before that, lets enjoy a **Gauss story**.

$$\sum_{i=1}^{100} i = 1 + 2 + 3 + \dots + 100 = ?$$

$$= (1+100) + (2+99) + (3+98) + \dots + (50 + 51)$$

$$= 101 + 101 + 101 + \dots + 101$$







1777 - 1855

<u>**Tip #1:</u>** Pulling out a final term from a summations</u>

- In working with summations, it is sometimes useful to be able to *pull out* (or *add in*) a *final term* to a summation.
- This is often done so we can use a specific closed formula (*e.g.*, $\frac{n(n+1)}{2}$)
- For *n* > *m*, we have the following trick:

$$\sum_{k=m}^{n} (a_k) = \sum_{k=m}^{n-1} (a_k) + a_n$$

Try the following:

Answers:

$$\sum_{j=0}^{n} (2j) = ??$$

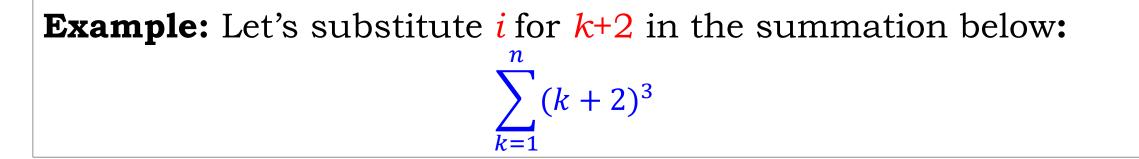
$$\sum_{j=0}^{n} (2j) = \sum_{j=0}^{n-1} (2j) + 2n$$

$$\sum_{j=0}^{n+2} \left(2^{j-1} \right) = ??$$

$$\sum_{j=0}^{n+2} (2^{j-1}) = \sum_{j=0}^{n+1} (2^{j-1}) + 2^{n+1}$$

<u>Tip #2</u>: Change of variables in summations

- The variable used for the index in a summation is internal to the sum and can be *replaced* with any other variable name (i.e., the value of *i*, *j*, *k*).
- Substitutions can be done for the index variable and will require that both the *upper and lower limit be adjusted*.

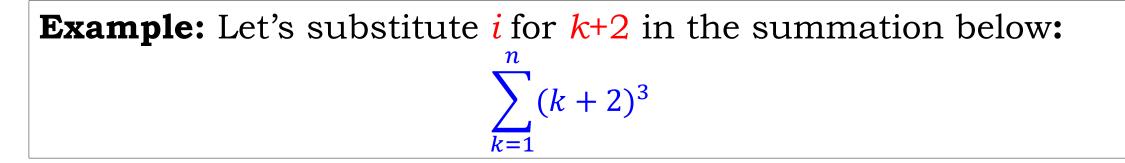


We want to have an expression of the form

What should be the correct upper and lower limits (in terms of i)?

 $\sum^{i} i^{3}$

Lower limit: When k = 1, what should be the value of i? We have selected, i = k + 2So, $k = 1 \implies i = 3$



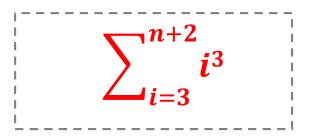
Similarly,

<u>Upper limit</u>: When k = n, what should be the value of *i*?

We have selected, i = k + 2

So, $k = n \implies i = n + 2$

So, we get



When analyzing an algorithm we often run across coding that involves some type of loop and we need to "count" the number of times the loop is executed.

Sometimes things are pretty *easy* to count ...

```
FOR i = 1 to n
    sum = sum + i //do this statement n times
End FOR
```

When analyzing an algorithm we often run across coding that involves some type of loop and we need to "count" the number of times the loop is executed.

Sometimes things are pretty <u>not so easy</u> to count ...

When it is not easy to count, we proceed by attempting to find a **closed form** to assist in counting operations.

A **closed form** is an expression that can be computed by applying a fixed number of familiar operations to arguments (e.g., n(n+1)).

Determining the closed form can make it easier than trying to count the number of operations in the following expression: 2 + 4 + ... + 2n (not closed form).

Some helpful facts for deriving closed forms:

1.
$$\sum_{k=m}^{n} c = c(n-m+1)$$
 [Sum of constant]

2. $\sum c a_k = c \sum a_k$ [Sum of constant]

3.
$$\sum_{k=1}^{n} (a_k - a_{k-1}) = a_n - a_0$$
 [Collapsing sum]

4.
$$\sum_{k=1}^{n} (a_{k-1} - a_k) = a_0 - a_n$$
 [Collapsing sum]

5.
$$\sum (a_k + b_k) = \sum a_k + \sum b_k$$
 [Sum of sums]

6.
$$\sum_{k=m}^{n} a_{k+i} = \sum_{\substack{k=m+i \ k=m+i}}^{n+i} a_k$$

[Shifting index]

7.
$$\sum (a_k x^{i+k}) = x^i \sum (a_k x^k)$$

8.
$$\sum_{k=1}^{n} k = \frac{n(n+1)}{2}$$

9.
$$\sum_{k=1}^{n} k^2 = \frac{n(n+1)(2n+1)}{6}$$

10. $\sum_{k=0}^{n} a^k = \frac{a^{n+1}-1}{a-1}; a \neq 1$

11.
$$\sum_{k=1}^{n} k a^k = \frac{a - (n+1)a^{n+1} + na^{n+2}}{(a-1)^2}; a \neq 1$$

Let's take a look at how we might convert a given summation into a closed form:

Example: Find a closed form for the following:

 $3 + 7 + 11 + \dots + (3 + 4n)$

Step 1: Get a handle on the summation = 3 + 7 + 11 + ... + (3 + 4n) = $\sum_{k=0}^{n} (3 + 4k)$ Step 2: Try to simplify or apply any known tricks = 3(n+1) + $\sum_{k=0}^{n} (4k) = 3(n+1) + \sum_{k=1}^{n} (4k)$ = 3(n+1) + 4 $\sum_{k=1}^{n} k = 3(n+1) + 4(\frac{n(n+1)}{2})$ = 3 + 3n + 2n² + 2n = $2n^2 + 5n + 3$

Try the following.

Example: Find a closed form for the following expression:

$$\sum_{k=1}^{n} (2k+3)$$

Answer:

$$\sum_{k=1}^{n} (2k+3) = 3n+2\sum_{k=1}^{n} k$$
$$= 3n+2\left(\frac{n(n+1)}{2}\right) = 3n+n(n+1) = \boxed{n^2+4n}$$

Example: Find a closed form for the following sequence: 4 + 8 + 12 + 16 + ... + 4n

[Try to find the pattern (find the *i*th term)]

[Write the summation expression]

[Simplify – Find the closed form]

$$= 4 \times \frac{n(n+1)}{2} = 2n^2 + 2n$$

4i

 $= 4 \sum_{i=1}^{n} i$

4*i*

Example: Find a closed form for the following (a bit challenging): $2 + (2^2 \times 7) + (2^3 \times 14) + (2^4 \times 21) + (2^n \times 7(n-1))$

Answer:

$$2^i \times 7(i-1)$$

[Try to find the pattern (find the *i*th term)]

$$= 2 + \sum_{i=1}^{n} 2^{i} \times 7(i-1)$$

[Write the summation expression]

= $2 + 7 \sum_{i=1}^{n} 2^{i}(i-1)$ [Simplify – Take constant out (Fact 2)]

Example: Find a closed form for the following (a bit challenging): $2 + (2^2 \times 7) + (2^3 \times 14) + (2^4 \times 21) + (2^n \times 7(n-1))$

Answer:

$$2^i \times 7(i-1)$$

[Try to find the pattern (find the *i*th term)]

$$= 2 + \sum_{i=1}^{n} 2^{i} \times 7(i-1)$$

[Write the summation expression]

$$= 2 + 7 \sum_{i=1}^{n} 2^{i}(i-1)$$

[Simplify – Take constant out (Fact 2)]

$$= 2 + 7\sum_{i=1}^{n} (2^{i}i - 2^{i}) = 2 + 7\left(\sum_{i=1}^{n} 2^{i}i - \sum_{i=1}^{n} 2^{i}\right)$$

[Simplify – (Fact 5)]

Example: Find a closed form for the following (a bit challenging):

$$2 + (2^2 \times 7) + (2^3 \times 14) + (2^4 \times 21) + (2^n \times 7(n-1))$$

Answer: (continued) = $2 + 7 \left(\sum_{i=1}^{n} 2^{i} i - \sum_{i=1}^{n} 2^{i} \right)$

$$= 2 + 7\left(\left(2 - (n+1)2^{n+1} + n2^{n+2}\right) - \sum_{i=1}^{n} 2^{i}\right)$$
 [Simplify - (Fact 11)]

$$= 2 + 7\left(\left(2 - (n+1)2^{n+1} + n2^{n+2}\right) - \left(2^{n+1} - 1 - 1\right)\right)$$

[Simplify - (Fact 10)

Example: Find a closed form for the following (a bit challenging):

$$2 + (2^2 \times 7) + (2^3 \times 14) + (2^4 \times 21) + (2^n \times 7(n-1))$$

Answer: (continued)

$$= 2 + 7\left(\left(2 - (n+1)2^{n+1} + n2^{n+2}\right) - \left(2^{n+1} - 1 - 1\right)\right)$$

$$= 2 + 7 \left(2 + (n-1)2^{n+1} - \left(2^{n+1} - 1 - 1 \right) \right)$$

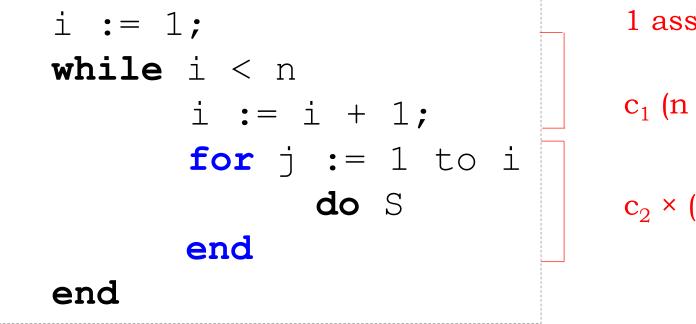
[Simplify green term]

$$= 2 + 7 \left(2 + (n-1)2^{n+1} - \left(2^{n+1} - 2 \right) \right)$$

 $= 2 + 7 \left(4 + (n-2)2^{n+1} \right)$

[Simplify red term]

Example: Let count(n) be the number of := statements executed by the following algorithm as a function of n. Find a closed form for a count(n).



1 assignment op $c_1 (n - 1)$ $c_2 \times (2 + 3 + ... + n)$ ops

Answer:

We ignore constants c_1 and c_2 .

The total count of operations, count(*n*) is:

$$= 1 + (n-1) + (2 + 3 + 4 + ... + n)$$

$$= 1 + 2 + 3 + 4 + \dots + n + (n - 1) = \frac{n(n+1)}{2} + (n - 1)$$

$$=\frac{n(n+1)}{2} + \frac{2(n-1)}{2} = \frac{n^2 + n + 2n - 2}{2} = \frac{n^2 + 3n - 2}{2}$$

Lets try a harder one.

Example: Let count(n) be the number of := statements executed by the following algorithm as a function of n. Find a closed form for a count(n).

(1)

i := 1;
while i < n do
 i := i + 2;
 for j := 1 to i
 do S
 end
end</pre>

(n-1)/2 (igoes up by 2) ???

Let's pick a specific n (i.e., n = 8) to just try to get a handle on what the heck is happening inside the loop.

| i := 1; |
|-------------------------------|
| while i < n do |
| i := i + 2; |
| for j := 1 to i |
| do S |
| end |
| end |

| i | i < 8? | j |
|---|--------|--------|
| 1 | Т | 1 to 3 |
| 3 | Т | 1 to 5 |
| 5 | Т | 1 to 7 |
| 7 | Т | 1 to 9 |
| 9 | F | |

Calls to S

3

5

7

9

2*k*+1

Let's generalize it now for any *n*.

i < n? i := 1; 1 to 3 Т while i < n do 1 i := i + 2;3 Τ 1 to 5 for j := 1 to i 5 1 to 7 Τ do S 7 Τ 1 to 9 end 2*k*-1 Τ 1 to 2k+1end 2*k*+1 F

So the number of calls to S looks like

Calls to S:

$$3 + 5 + 7 + \dots + (2k + 1)$$

= $\sum_{h=1}^{k} (2h + 1)$ = $\sum_{h=1}^{k} (2h) + \sum_{h=1}^{k} (1)$
= $2 \times \frac{k(k+1)}{2} + k$
= $k^2 + 2k$

We now know the total number of calls to S. But what is $k^2 + 2k$ in terms of *n*? That's really what we need to know.

- Question on the table: What is $(k^2 + 2k)$ in terms of *n*?
- We can plug that result back into the closed form on the previous slide for *k* to get the total count of calls to S.

| i | i < n? | j |
|---------------|--------|--------------------|
| 1 | Т | 1 to 3 |
| 3 | Т | 1 to 5 |
| 5 | Т | 1 to 7 |
| 2 <i>k</i> -1 | Т | 1 to 2 <i>k</i> +1 |
| 2 <i>k</i> +1 | F | |

| $(2k-1) < n \leq (2k+1)$ |
|-------------------------------|
| $(2k-2) < (n-1) \leq (2k)$ |
| $(k-1) < \frac{n-1}{2} \le k$ |

Summary:

