

# Improving Security of the Wireless Sensor Networks by Implementing a Secure Routing Scheme

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**Abstract-** A Wireless Sensor Network is a self-configuring network of small sensor nodes communicate amongst themselves using radio alerts, and deployed in amount to experience, display and understand the physical global. Designing comfortable authentication mechanisms in wi-fi sensor networks with the intention to companion a node to a at ease network is not an clean undertaking because of the constraints of this form of networks. While data sharing, the sender node must verify it's neighbor nodes either genuine or not. If the neighbor node may have malicious behavior then our data will be collapsed. In this paper we describing concept to provide security to wireless sensor networks by distributing keys to all nodes in a network and whenever any source transfer data then it will encrypt data and transfer to destination via neighbors and if neighbors have proper keys can only decrypt and reencrypt and send data to destination.

**Keywords-** *Wireless Sensor Networks Routing; Malicious Nodes; Neighbor Nodes; Encryption and Decryption*

## I. INTRODUCTION

Wireless Sensor Networks (WSNs) can be defined as a self-configured and infrastructureless wi-fi networks to monitor physical or environmental situations, together with temperature, sound, vibration, stress, movement or pollutants and to cooperatively skip their facts through the network to a primary region or sink in which the data can be discovered and analyzed. A sink or base station acts like an interface between customers and the network. One can retrieve required statistics from the network by way of injecting queries and collecting outcomes from the sink. Typically a wi-fi sensor network includes masses of thousands of sensor nodes. The sensor nodes can communicate among themselves using radio alerts. A wi-fi sensor node is ready with sensing and computing devices, radio transceivers and power components. The individual nodes in a wi-fi sensor network (WSN) are inherently resource limited: they've limited processing speed, garage potential, and conversation bandwidth. After the sensor nodes are deployed, they're responsible for self-organizing the

appropriate network infrastructure often with multi-hop network with them. Then the onboard sensors begin gathering facts of hobby. Wireless sensor gadgets additionally reply to queries despatched from a "manipulate website" to carry out precise instructions or provide sensing samples. The working mode of the sensor nodes may be either non-stop or event driven. Global Positioning System (GPS) and local positioning algorithms can be used to achieve area and positioning information. Wireless sensor gadgets may be equipped with actuators to "act" upon sure conditions. These networks are from time to time greater in particular referred as Wireless Sensor and Actuator Networks as described. Wireless sensor networks (WSNs) permit new programs and require non-conventional paradigms for protocol layout due to numerous constraints. Owing to the requirement for low device complexity together with low electricity intake (i.e. Long network lifetime), a proper balance between communication and sign/facts processing capabilities ought to be determined. This motivates a massive effort in studies sports, standardization system, and modern ventures on this field since the most recent decade. At exhibit time, the greater part of the examination on WSNs has focused on the outline of vitality and computationally productive calculations and conventions, and the application space has been limited to basic information arranged observing and revealing applications. The creators propose a Cable Mode Transition (CMT) calculation, which decides the negligible number of dynamic sensors to keep up K-scope of a landscape and K-availability of the system. In particular, it allots times of dormancy for link sensors without influencing the scope and availability prerequisites of the system construct just in light of nearby data. A postponement mindful information accumulation arranges structure for remote sensor systems is proposed. The target of the proposed organize structure is to limit delays in the information gathering procedures of remote sensor systems which expands the lifetime of the system. The writers have considered transfer hubs to alleviate the system geometric inadequacies and utilized Particle Swarm Optimization (PSO) based calculations to find the ideal sink area regarding those hand-off hubs to conquer the lifetime challenge. Vitality proficient correspondence has additionally

been tended. The creators proposed a geometrical answer for finding the ideal sink situation for expanding the system lifetime. More often than not, the explorations on remote sensor systems have thought about homogeneous sensor hubs. In any case, these days' analysts have concentrated on heterogeneous sensor systems where the sensor hubs are dissimilar to each other as far as their vitality. Few authors tend to the issue of sending hand-off hubs to furnish adaptation to internal failure with higher system availability in heterogeneous remote sensor systems, where sensor hubs have distinctive transmission radii. New system designs with heterogeneous gadgets and the ongoing headway in this innovation dispose of the present confinements and extend the range of conceivable applications for WSNs impressively and all these are changing quickly.

## II. RELATED WORK

S Zhu, S Setia, S Jajodia propose LEAP (Localized Encryption and Authentication Protocol); a key administration convention planned to help a few correspondence designs. In this convention, every hub stores four sorts of keys: individual, pairwise, bunch, and gathering. An individual key is a key shared between a hub and the base station. A pairwise key is shared between a hub and every one of its neighbors. A group key is a key shared between a hub and every single neighboring hub. A gathering key is a key normal to the whole system. The individual key is preloaded. After organization, neighboring hubs set up pairwise keys. They validate themselves utilizing a pre-conveyed key which is deleted when pairwise keys are built up. To build up bunch keys and the gathering key, hubs utilize communicates and message handing-off. The convention utilizes  $\mu$ Tesla to validate communicates. B Lai, S Kim, I Verbauwhede propose BROSK (BROadcast Session Key arrangement convention). With BROSK each hub communicates a message containing its nonce. Thus, every two neighboring hubs that hear each other can figure a typical key which is an element of their two nonces. Neighboring hubs confirm themselves with a predeployed key which should be inaccessible for the situation the hub is caught. As of late Han et al. propose verification demonstrate that goes for lessening overhead for the re-confirmation of sensor hubs utilizing symmetric and open key cryptography. It depends on a ticket scrambled utilizing a typical mystery key between neighboring settled hubs. This ticket is sent to a portable hub amid the primary validation stage. This ticket is just helpful when the versatile hub chooses to re-verify with this neighbor settled hub. What's more, the convention just functions admirably when the settled hub is in coordinate range with the base station, and the underlying confirmation stage experiences interior assaults. Sinks in the system can without much of a stretch replace each other when they are not in correspondence go with the base station. The creators just give a snappy casual security

investigation of their answers and they have not sent their answers on genuine bits.

Liu et al. propose LBKs (area based keys) that depends on area data to accomplish key administration. The keys are set up as indicated by the land area of sensor hubs. In any case, knowing the geological area of hubs isn't ensured with arbitrary organization. Eschenauer and Gligor propose a plan in light of an arbitrary key pre-dispersion. In this plan, every sensor arbitrarily picks an arrangement of keys and their identifiers from a key pool before sending. At that point, a mutual key revelation stage is propelled where two neighbors trade and think about rundown of personalities of keys in their key chains. Fundamentally, every sensor hub communicates one message and gets one message from every hub inside its radio range where messages convey key ID records. In this way, any match of hubs has a specific likelihood to share no less than one regular key. The test of this plan is to locate a decent exchange off between the measure of the key pool and the quantity of keys put away by hubs to accomplish the best likelihood. The principle disadvantage of this approach is that if the quantity of traded off hubs builds, the division of influenced connects likewise increments. R Anderson, H Chan, A Perrig center around creating cost-sparing instruments while debilitating the risk display; they propose Key Infection, a lightweight security convention appropriate for use in noncritical item sensor systems where an aggressor can screen just a settled level of correspondence channels.

## III. FRAMEWORK

### A. Proposed System

In this paper author describing concept to provide security to wireless sensor networks by distributing keys to all nodes in a network and whenever any source transfer data then it will encrypt data and transfer to destination via neighbors and if neighbors have proper keys can only decrypt and re-encrypt and send data to destination. If any attacker node arrives by stealing neighbor's id then cant able to identify key and can't decrypt data to do malicious activity. So, by adding key distributing mechanism and data cryptography with keys can avoid attacker node from doing malicious activity. To implement above concept we did modification to DSR protocol to support keys and cryptography techniques and for extension work we did modification to AODV protocol to support same activity.

### B. Dynamic Source Routing Protocol (DSR)

The distinguishing features of DSR are: low network overhead, no extra infrastructure for management and the use of source routing. Source routing means that the sender had complete understanding of the entire hop-with the aid of-hop direction information to the vacation spot. The protocol consists of the 2 essential mechanisms of Route Discovery

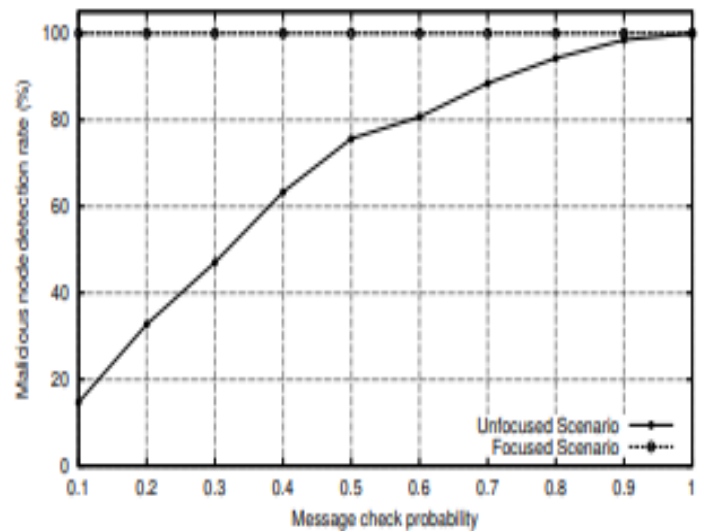
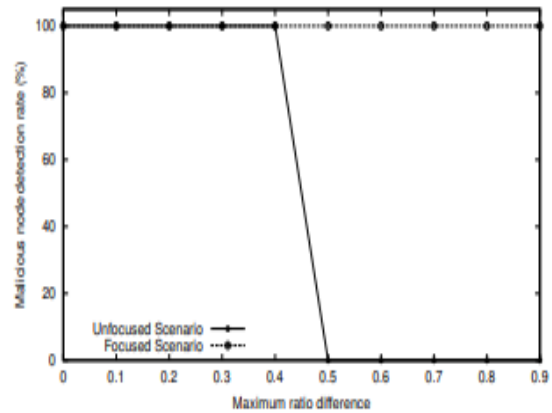
and Route Maintenance. Normally routes are stored in a path cache of each node. When a node loves to speak to a vacation spot, first it exams for the course for that precise vacation spot in the direction cache. If sure, the packets are sent with supply course header information to the vacation spot. In the other case, if the direction is not to be had on the direction cache; then the node will initiate the direction discovery mechanism to get the route first. The direction discovery mechanism will flood the network with route request (RREQ) packets, after which the buddies will get hold of RREQ packets and test for the path to destination of their direction cache. If the course isn't in their caches rebroadcast the RREQ, otherwise the node replies to the originator with a route error (RREP) packet. Since RREQ and RREP packets each are source routed, original source can reap the course and upload to its route cache. In any case the hyperlink on a source route is damaged; the source node is notified with a path errors (RERR) packet. Once the RERR is acquired, the supply gets rid of the route from its cache and direction discovery method is reinitiated. DSR being a reactive routing protocol has no need to periodically flood the network for updating the routing tables like table-pushed routing protocols do. Intermediate nodes are capable to make use of the path cache statistics efficaciously to lessen the manage overhead.

C. Neighbor-Based Malicious Node Detection

In detecting malicious nodes in the presence of faults and occasions, we appoint a smoothing filter and self belief stage assessment to enhance the malicious node detection fee. A filter out is used to accurate a few false readings because of transient faults. It as a result efficaciously reduces the temporary fault possibility pt in any such way that malicious nodes may be detected for a wider range of malicious node chance. Confidence ranges are hired to estimate the trustworthiness of sensor nodes, reflect the tiers in decision making system, and logically isolate malicious nodes and nodes with a permanent fault from the network.

IV. EXPERIMENTAL RESULTS

In these experiments, nodes of the network do now not check all transmissions they listen. Instead, they do so with a sure probability, given by means of the test chance.



This indicates that it isn't always essential to test all transmissions one hears with a purpose to have a very good malicious node detection rate. On the opposite hand, inside the focused case, the detection rate is 1 for any take a look at possibility.

V. CONCLUSION

In this paper we described that in the wireless sensor networks are facing various problems in current scenario. In that the main problem is to detecting the malicious behavior node from the current routing path. In this paper we performed few modifications in the DSR protocol and it can detect the malicious behavior nodes in the routing path. By implementing the proposed scheme, we can enhance the performance of the wireless sensor routing protocols.

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