

# Secured Data Sharing in Group Users in Cloud

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**Abstract-** Data security becomes essential every day with development and growth of cloud computing. Cloud computing enables distributed resources and services to be shared between different organisations or sites. Ex: Microsoft Azure, AWS. Since group sharing and distributed resources in an open environment through network which rises security problem. By SeDaSC (Secure Data Sharing in Clouds) methodology, we can provide, data sharing securely among group users without re-encryption and zero proof of knowledge is used to provide trust level at Cryptographic Server side.

**Keywords-** Access control, Cloud computing, Confidentiality, Secure Sharing, Security, Trusting.

## I. INTRODUCTION

“Cloud computing is quickly emerging due to the provisioning of versatile, adaptable, and on-request storage and computing administrations for clients. Associations with a low budget would now be able to use high computing and storage services without vigorously investing in infrastructure and maintenance. In addition, the secrecy the executives by a client guarantees that the cloud does not take in any information about the client data. Cryptography is utilized as a common instrument to give confidentiality and protection administrations to the data. The data are usually encrypted before storing to the cloud”.

The data handling of a group has certain additional qualities in contrast to the two-party correspondence or data handling of a single customer. “[1] the existing, departing, and recently joining group individuals can turn out to be an insider danger violating data confidentiality and protection. Insider dangers can demonstrate to be more devastating because of the way that they are for the most part launched by confided in elements. Because of the way that individuals trust insider identities, the examination network concentrates more on outsider attackers. In any case, various security issues can emerge due to distinctive users in a group. We examine a portion of the issues in the following discussion. A single key shared between all group individuals will result in the access of past data to a recently joining part. The foresaid situation disregards the privacy and the principle of least privilege”.

“SeDaSC that bargains with the forementioned security necessities of shared group data within the cloud. The SeDaSC approach works with three elements as follows: users, a cryptographic server and cloud. The data proprietor presents

the data, the rundown of the users, and the parameters required for generating an entrance control list (ACL) to the CS. The CS is a confided in outsider and is mindful for key the executives, encryption, unscrambling, and get to control. The CS produces the symmetric key and scrambles the data with the created key”.

## II. LITERATURE SURVEY

[1] *Khan and Abbas*, “The framework is structured with the goal that it will catch the state of patient at all stages. There is no compelling reason to find the patient's past restorative records volume however it ensures the exactness in data and legitimate drug to the patient. But EHR integration i.e. the procedure of patient information sharing among medicinal services suppliers and exchanging them over internet with other human services supplier remains a challenge and genuine worry since it is exposing to robbery, security violation. The framework which was actualized in country area of China says around individuals were included and aggregated using this EHR system. The storage of the wellbeing records needs the infrastructure which guarantees the secured storage and accessibility at whenever. The movement of the wellbeing records from paper to these paperless e-Records has a more noteworthy alleviation among the service providers and furthermore for the doctor's facilities. This framework likewise provides the more noteworthy correspondence between the city healing facilities and the rural clinics”.

[3] *Madani, A. N. Khan, M. L. M. Kiah*, “The operating framework oversees the computer equipment assets and gives an interface to interaction of client and application programming with equipment resources. The hypervisor is framework programming that enables users to remotely create virtual machines on cloud server(s) at runtime. The virtual machine has client defined equipment specifications and a software stack. The virtualization procedure enhances the accessibility of the user's facilitated administrations even in the event of equipment disappointment. The virtual machine with the whole programming stack can be relocated to another server with insignificant inaccessibility of facilitated administrations. Moreover, virtualization is additionally helpful for cloud specialist organizations. The dispersed physical servers without virtualization use just 20% of aggregate capabilities. The virtualization process can support equipment usage. The middleware framework programming manages the straightforward execution and interaction among employments running on cloud servers. The product

infrastructure layer hands over the organized assets to upper layers and gives a foundation for another computing worldview that conveys IT as an administration”.

[4] X. Wu and Zhang, “A average methodology for data privacy security is to encrypt a data with a (normally symmetric) key before storing it to cloud. In any case, encryption makes it troublesome for adaptable sharing data between various users. On one side, sharing the data encryption key to all users effortlessly empowers a user to get to all data that put away in cloud of a data owner, which abuses the minimum benefit principle. This introduces overwhelming computing load at data proprietor side, and relies on trust and key administration frameworks, for example, PKI which may not be scale and adaptable well. Communicate encryption and group key administration can be utilized for sharing data in group way. In any case, managing group is complicated in specific for the present inescapable data sharing such as cloud-based joint efforts and interpersonal organizations, where the number of groups of interests for an individual client is large. In expansion, the span of a group can likewise be vast, and the membership for the most part changes much of the time, which makes group key the board exceptionally repetitive for a typical client”.

### III. PROBLEM STATEMENT

Here unique key shared among all the group individuals will result in the entrance of past data to a recently joining part. “The aforesaid circumstance disregards the privacy and the principle of slightest benefit. In like manner, a departing part can get to future correspondence. Therefore, in group-shared data, the inside individuals may create the issue of in reverse access control (another client accessing past data) and forward access control (a departing client accessing future data). The straightforward arrangement of rekeying (generating another key, decrypting every one of the data and re-encrypting with the new key) does not turn out to be scalable for frequent changes in the group enrollment”.

A certificate less proxy re-encryption (CL-PRE) plot for securely sharing the data within a group in people in general cloud. “In the CL-PRE plot, the data proprietor encodes the data with the symmetric key. Along with the symmetric key is scrambled with people in general key of the data proprietor. Both the scrambled data and the key are transferred to the cloud. The public-private keys created in the proposed plan are not founded on the certificates. The client's character is utilized to create the public private key combine. The proxy re-encryption depends on bilinear pairing and the BDH that makes the CL-PRE plot computationally intensive. The computational expense of the bilinear pairing is high as compared with the standard tasks in finite fields”.

### IV. PROPOSED METHOD IMPLEMENTATION

“The SeDaSC technique works with mainly three substances like Users, Cryptographic server and Cloud. The data proprietor presents the data, the rundown of the users, and the parameters required for generating an ACL to the CS. The CS is a confided in outsider and is responsible for key administration, encryption, decoding, and access control. The CS creates the symmetric key and scrambles the data with the produced key”.

Along these lines, “for every client in the group, the CS isolates the key into two sections to such an extent that a single part alone can't regenerate the key. Progressively, the original key is erased through secure overwriting. The scrambled data are hence transferred to the cloud for storage for the benefit of the client. The client who wishes to get to the data sends a download request to the CS. The CS, in the wake of authenticating the requesting client, receives the bit of the key from the client and in this manner downloads the data record from the cloud. The data are decoded and sent back to the client. For a recently joining part, the two bits of the key are created, and the client is added to the ACL. For a departing part, the record is erased from the ACL. The departing part can't unscramble the data all alone as he/she just has a bit of the key. Also, no frequent unscrambling and re-encryption are required in the event of changes in the group participation”.

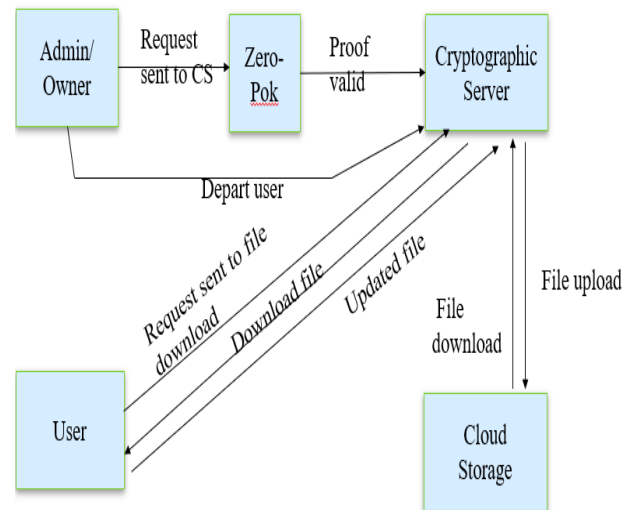


Fig. 1: Proposed method system work flow

The methodology of the SeDaSC has the following entities.

#### a. Cloud

“The cloud gives storage administrations to the client. The data on the cloud should be secured against protection breaches. The secrecy of the data is ensured by storing encrypted data over the cloud. The cloud in the SeDaSC

methodology only involves essential cloud tasks of record transfer and download. Therefore, no progressions at the convention or implementation level on the cloud are required”.

#### b. Cryptographic Server

“The CS is a confided in gathering and is responsible for security operations, for example key management, encryption, decoding, the management of the ACL for providing privacy, and secured data forwarding among the group. The users of SeDaSC are required to be registered with the CS to obtain the security services. The CS is thought to be a secure element in the proposed methodology. The CS can be maintained by an association or can be claimed by an outsider supplier”. Be that as it may, the cryptographic server maintained by an association will create additional trust in the framework.

#### c. Users

“The users are the customers of the storage cloud for each data record; one client will be the owner of the document, whereas the others in the group will be the data shoppers. The record chooses the entrance privileges of the other group members. The get to rights are conceded and revoked dependent on the decision of the proprietor. The entrance rights are overseen by the CS in the form of an ACL record. A different ACL is maintained for each of the data records”.

#### d. Cryptographic Keys

“The SeDaSC technique maintains a single cryptographic key for every one of the data documents. In any case, after encryption/decryption, the entire key isn't stored and controlled by any of the involved gatherings. The key is apportioned into two constituent parts and are controlled by different elements. The following are the keys that are utilized within SeDaSC”.

#### e. Symmetric Key as K

K is an irregular secret created by the CS for every one of the data documents. “The length of K in SeDaSC is 256 bits, as is recommended by the greater part of the norms regarding key length for symmetric key calculations (SKAs). In any case, the length of the key can be altered according to the requirements of the underlying SKA. K is obtained in a two-advance process”.

In the initial step, “an arbitrary number R of length 256 bits is generated with the end goal that  $R = \{0, 1\}^{256}$ . In the subsequent stage, R is passed through a hash work that could be any hash work with a 256-bit yield. For our situation, we utilized secure hash algorithm 256 (SHA-256). The second step totally randomizes the initial client determined irregular number R. The yield of the hash function is named as K and is utilized in symmetric key encryption [e.g., the Advanced Encryption Standard (AES)] for securing the data”.

#### f. CS Key Share as $K_i$

For all users of the group, the CS creates  $K_i$ , with the end goal that  $K_i = \{0, 1\}^{256}$ . “ $K_i$  fills in as the CS segment of the key and is utilized to process K at whatever point an encryption/decoding request is received by the CS.

Moreover, it is ensured by correlation that the distinct  $K_i$  is created for every record client”.

#### g. User Key Share as $K_i$

$K_i$  is calculated for each group of users by XOR of K and  $K_i$ .

#### h. Zero Proof of Knowledge (Zero PoK)

In our framework data proprietor sending a record to CS for encryption the document. Be that as it may, sending a record to CS directly isn't protected, so to utilize Zero PoK, we can know whether our document is sending legitimate CS or not.

## V. ALGORITHM

In cryptography, let's say that Alice wants to prove that Bob knows a value (x), such that:

$$g^x \bmod P = Y$$

Where g is a pre-selected value, P is a prime and Y is a result.

Both Bob and Alice know these values, and it's difficult to know the value of x, as there are many values of x that would fit.

To prove that Bob knows the value of x, he creates a random number (r) and sends the result of this calculation to Alice:

$$C = g^r \bmod P$$

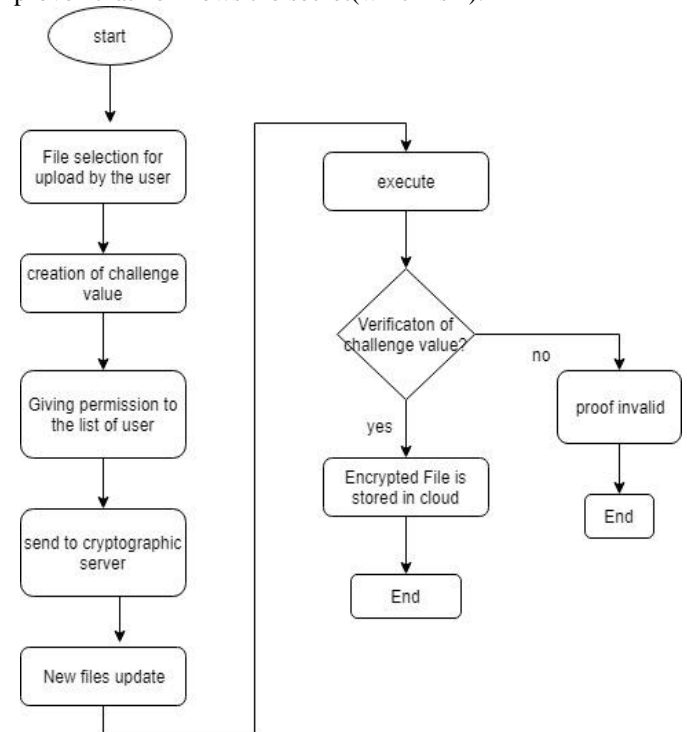
He then sends:

$$\text{Cipher1} = g^{(x+r) \bmod (P-1)} \bmod P$$

Alice then calculates:

$$\text{Cipher2} = C \cdot Y \bmod P$$

If the values are the same (Cipher1 equals Cipher2), Bob has proven that he knows the secret (which is x).



#### IV. PERFORMANCE RESULTS

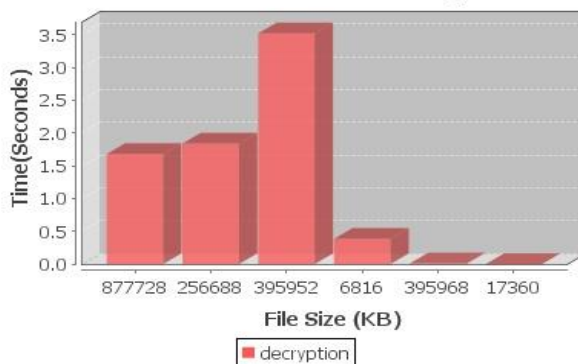
We additionally assessed the SeDaSC methodology based on the aggregate time devoured to transferring/downloading a record to/from the cloud. “The aggregate time is composed of the time from the season of accommodation of request to the CS to the point of time at which the document is transferred/downloaded to/from the cloud. The following occasions are included in the total time”:

“key computation time, encryption/decryption time, upload/download time. The time of request and other related data submission to the CS and the cloud”.

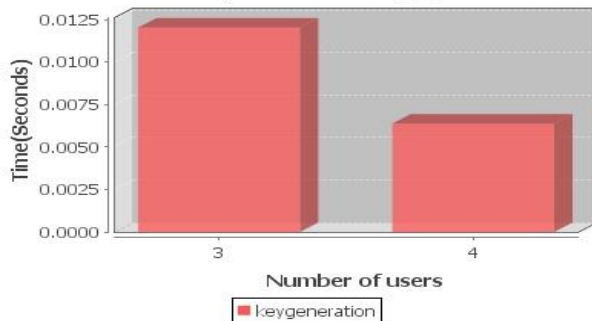
##### Performance of File Encryption



##### Performance of File Decryption



##### Time consumption for key generation



#### V. CONCLUSION

Cloud computing offers people great convenience. In particular, it corresponds perfectly to the increased need for the Internet to share data. Building a cost-efficient and secure cloud data sharing system, the SeDaSC methodology provides secured data sharing between group users without re-encryption for each and every user. For increasing the trust level at Cryptographic side, using Zero proof of knowledge.

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